



UNIVERSITY OF MINES AND TECHNOLOGY, TARKWA

FIRST SEMESTER EXAMINATIONS, NOV/DEC 2018

COURSE NO: CE 379

COURSE NAME: INFORMATION THEORY

CLASS: CE III

TIME: 3 HOURS

Name: _____ Index Number: _____

Answer ALL questions in Section A and any other TWO in Section B. Circle your answers in Section A.

SECTION A

1. Two cards are drawn at random from a pack of 52 cards. The probability that both are spades is
 - a. $1/15$
 - b. $2/17$
 - c. $1/17$
 - d. $2/15$
2. Following is not a unit of information
 - a. Hz
 - b. nat
 - c. bit
 - d. decit
3. The channel capacity of a noise free channel having M symbols is given by
 - a. $\log M$
 - b. 2
 - c. M
 - d. None
4. Information content in a universally true event is
 - a. Infinite
 - b. Positive Constant
 - c. Negative Constant
 - d. Zero
5. As the bandwidth approaches infinity, the channel capacity becomes
 - a. $1.44S/\eta$
 - b. Infinite
 - c. Zero
 - d. None of the above
6. Entropy represents
 - a. Information Signal
 - b. Amplitude of Information
 - c. Average Information per Message
 - d. All of the above
7. The ideal communication channel is defined for a system which has
 - a. finite C
 - b. $BW=0$
 - c. $S/N = 0$
 - d. Infinite C
8. The channel capacity is a measure
 - a. entropy rate
 - b. maximum rate of information a channel can handle
 - c. information contents of messages transmitted in a channel
 - d. None of the above
9. The probability of a message is $1/16$. The information in bits is
 - a. 1 bit
 - b. 2 bits
 - c. 3 bits
 - d. 4 bits

10. Three light bulbs are chosen at random from 15 bulbs out of which 5 are defective. The probability that exactly one is defective
- $24/91$
 - $45/91$
 - $67/91$
 - None of the above
11. The probability density function of a Markov process is
- $p(x_1, x_2, x_3, \dots, x_n) = p(x_1)p(x_2/x_1)p(x_3/x_2) \dots p(x_n/x_{n-1})$
 - $p(x_1, x_2, x_3, \dots, x_n) = p(x_1)p(x_1/x_2)p(x_2/x_3) \dots p(x_{n-1}/x_n)$
 - $p(x_1, x_2, x_3, \dots, x_n) = p(x_1)p(x_2)p(x_3) \dots p(x_n)$
 - $p(x_1, x_2, x_3, \dots, x_n) = p(x_1)p(x_2 * x_1)p(x_3 * x_2) \dots p(x_n * x_{n-1})$
12. For M equally likely messages, the average amount of information H is
- $H = \log_{10}M$
 - $H = \log_2M$
 - $H = \log_{10}M^2$
 - $H = 2\log_{10}M$
13. Applications that require image compression include
- Internet
 - Multimedia
 - Medical Imaging
 - All of the above
14. Information Rate
- Information per Unit Time
 - Average Number of bits of Information Per Second
 - rH
 - All of the above
15. For M equally likely messages, $M > 1$, if the rate of information $R \leq C$, the probability of error is
- Arbitrarily small
 - Close to unity
 - Not predictable
 - Unknown
16. For a binary symmetric channel, the random bits are given as
- Logic 1 given by probability P and logic 0 by (1-P)
 - Logic 1 given by probability 1-P and logic 0 by P
 - Logic 1 given by probability P^2 and logic 0 by 1-P
 - Logic 1 given by probability P and logic 0 by $(1-P)^2$
17. The mutual information
- Is symmetric
 - Always non negative
 - Both a and b are correct
 - None of the above
18. The relation between entropy and mutual information is
- $I(X;Y) = H(X) - H(X/Y)$
 - $I(X;Y) = H(X/Y) - H(Y/X)$
 - $I(X;Y) = H(X) - H(Y)$
 - $I(X;Y) = H(Y) - H(X)$
19. The memory less source refers to
- No previous information
 - No message storage
 - Emitted message is independent of previous message
 - None of the above

20. The information I contained in a message with probability of occurrence is given by (k is constant)
- a. $I = k \log_{21} P$
 - b. $I = k \log_2 P$
 - c. $I = k \log_{21} 2P$
 - d. $I = k \log_{21} P^2$
21. When the base of the logarithm is 2, then the unit of measure of information is
- a. Bits
 - b. Bytes
 - c. Nats
 - d. none of the above
22. The probability of a 4 turning up at least once in two tosses of a fair die is
- a. $11/36$
 - b. $17/36$
 - c. $13/36$
 - d. $19/36$
23. The entropy for a fair coin toss is exactly
- a. 5 bits
 - b. 3 bits
 - c. 2 bits
 - d. 1 bit
24. Huffman Code produces a code which uses a minimum number of bits to represent each symbol
- a. True
 - b. False
25. For which value(s) of p is the binary entropy function $H(p)$ maximized?
- a. 0
 - b. 0.5
 - c. 1
 - d. 2
26. The more one learns, the more the uncertainty is
- a. True
 - b. False
27. The followings are the applicable areas of Information Theory
- a. Data Compression
 - b. Cryptography
 - c. DSL Modem
 - d. All of the above
28. The order of Commutation System is:
Source to Encoder to Channel to Decoder to Destination
- a. True
 - b. False
29. The major primary concern of Information Theory
- a. Source and Destination
 - b. Channel
 - c. Encoder and Decoder
 - d. All of the above
30. The long quest of finding computationally-efficient codes include:
- a. Lossless Source Coding (Huffman Code)
 - b. Noisy Channel Coding (Polar Code)
 - c. Lossy Channel Coding (Polar Code)
 - d. All of the above

SECTION B

1. Consider two random variables $X_1, X_2 \in \{0, 1\}$ with joint p.m.f

(x_1, x_2)	(0, 0)	(0, 1)	(1, 0)	(1, 1)
p(x ₁ , x ₂)	$\frac{1}{6}$	$\frac{1}{3}$	$\frac{1}{3}$	$\frac{1}{6}$

- a. Compute $H(X_1)$, $H(X_2)$, and $H(X_1, X_2)$
 - b. Compute $H(X_1|X_2 = 0)$, $H(X_1|X_2 = 1)$, $H(X_1|X_2)$, and $H(X_2|X_1)$.
 - c. Compute $I(X_1; X_2)$.
- 2 a. What do you understand by Self Information? Hence, state his properties
- b. *A source puts out one of five possible messages during each message interval. The probability of these messages are $\{m_1, \dots, m_5\} : P_1=1/2, P_2=1/4, P_3=1/4, P_4=1/16$ and $P_5=1/16$ What is the information content of these messages in bit?*
- c. Define Entropy. What are its mathematical properties?
- 2 a. What do you understand by channel capacity? Describe the developmental process of data compression algorithms.
- b. What are the application areas of Information Theory?

c. Consider the table below:

Symbol	P(S)
A	0.1
B	0.2
C	0.4
D	0.2
E	0.1

- i. Compute Entropy (H)
- ii. Build Huffman Code
- iii. Compute Average Code Length
- iv. Code "BCCADE"